

From linear to nonlinear n-width : optimality in reduced modelling

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Announcing two great conferences, hopefully in physical format

1. **Curves and Surfaces** : Arcachon, France, June 20-24, 2022

Registration open : cs2022.sciencesconf.org

2. **Foundation of Computational Mathematics** : Whistler, Canada, June 12-21, 2023.

Registration opening should be officially announced in coming months.

Save the dates !

Agenda

1. Linear widths
2. Widths of parametrized PDEs and reduced bases
3. Nonlinear widths
4. Widths and sampling numbers

Kolmogorov linear n -width

We are interested in approximating general functions $u \in V$, where V is a Banach space, by simpler functions v picked from a linear subspace $V_n \subset V$ of finite dimension n .

Classical Banach spaces : Lebesgue $L^p(\Omega)$, Sobolev $W^{m,p}(\Omega)$ for $\Omega \subset \mathbb{R}^d$.

Classical linear subspaces : algebraic or trigonometric polynomials of some prescribed degree, splines or finite elements on some given mesh, span of the n first elements $\{e_1, \dots, e_n\}$ from a given basis $(e_k)_{k \geq 1}$ of V .

Model class reflecting the properties the target function : $u \in \mathcal{K}$, where \mathcal{K} is a compact set of V . In parametrized PDEs, the set \mathcal{K} is the **solution manifold** that gathers all solutions $u(y) \in V$ as the parameter vector y varies.

Best choice of approximation spaces for this model class ?

The space V_n approximate \mathcal{K} with uniform accuracy

$$\text{dist}(\mathcal{K}, V_n)_V := \max_{u \in \mathcal{K}} \min_{v \in V_n} \|u - v\|_V$$

A.N. Kolmogorov (1936) defines the linear n -width of \mathcal{K} in the metric V as

$$d_n(\mathcal{K})_V := \inf_{\dim(V_n)=n} \text{dist}(\mathcal{K}, V_n)_V$$

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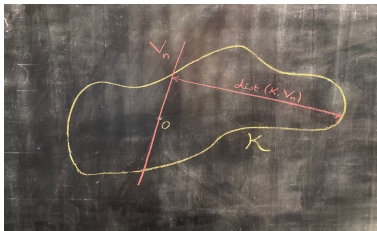
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Intuition



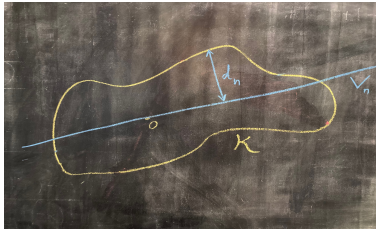
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may not exist. One often assumes it exists in order to avoid limiting arguments.

The quantity $d_n(\mathcal{K})_V$ can be viewed as a **benchmark/bottleneck** for numerical methods applied to the elements from \mathcal{K} that create approximations from linear spaces : interpolation, projection, least squares, Galerkin methods for solving PDEs...

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An analog concept in the stochastic framework : PCA

Assume that V is a Hilbert space and u is a random variable taking its value in V .

Optimal spaces in the mean-square sense.

$$\kappa_n^2 = \kappa_n(u)_V^2 := \min_{\dim(V_n)=n} \mathbb{E}(\|u - P_{V_n} u\|_V^2).$$

The space achieving the minimum is easily characterized by **principal component analysis** : consider the covariance operator

$$v \mapsto Rv = \mathbb{E}(\langle u, v \rangle u),$$

which is compact, when assuming that $\mathbb{E}(\|u\|_V^2) < \infty$. Diagonalized in the Karhunen-Loeve basis $(\varphi_k)_{k \geq 1}$ with eigenvalues $\lambda_1 \geq \lambda_2 \geq \dots \rightarrow 0$.

Then $V_n := \text{span}\{\varphi_1, \dots, \varphi_n\}$ and $\kappa_n^2 = \sum_{k > n} \lambda_k$.

Note that $\kappa_n(u)_V^2 \leq d_n(\mathcal{K})_V^2$ when u is supported in \mathcal{K} .

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Variants to n -width : realization of the approximation

The best approximation $u_n = \operatorname{argmin}\{\|u - v\|_V : v \in V_n\}$ is the orthogonal projection if V is a Hilbert space.

For a general Banach space, the map $u \mapsto u_n$ is not linear, and may even not be continuous (non-uniqueness of best approximation).

This motivates alternate definitions of widths where we impose linearity or continuity of the approximation process.

Approximation numbers are defined as

$$a_n(\mathcal{K})_V := \inf_L \max_{u \in \mathcal{K}} \|u - Lu\|_V,$$

with infimum taken over all **linear** maps L such that $\operatorname{rank}(L) \leq n$.

In a general Banach space $d_n \leq a_n \leq \sqrt{n}d_n$ and right equality may hold.

On the other hand one can prove that

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Behaviour of n -widths of smoothness classes

Typical compact sets in $V = L^p(\Omega)$ are balls of smoothness spaces. The behaviour of n -width is well understood for such sets. Example : $V = L^\infty(I)$ where $I = [0, 1] \subset \mathbb{R}$

and

$$\mathcal{K} = \mathcal{U}(\text{Lip}(I)) = \{u : \max\{\|u\|_{L^\infty}, \|u'\|_{L^\infty}\} \leq 1\},$$

Then one can prove

$$d_n(\mathcal{K})_V = \frac{1}{2n}, \quad n \geq 1, .$$

More generally when $V = W^{t,p}(\Omega)$ for some bounded Lipschitz domain $\Omega \subset \mathbb{R}^d$ and \mathcal{K} is the unit ball of $W^{s,p}(\Omega)$ with $s > t$, one can prove

$$cn^{-(s-t)/d} \leq d_n(\mathcal{K})_V \leq Cn^{-(s-t)/d}, \quad n \geq 1.$$

Curse of dimensionality : exponential growth in d of the needed n to reach accuracy ε .

Proof of upper bound : use a standard approximation method (piecewise polynomials, finite elements, or splines, on uniform partitions of Ω)

Proof of lower bound ? Two systematic approaches.

Bernstein width

Lemma : let $B_W = \{u \in W : \|u\|_V \leq 1\}$ be the unit ball of a subspace $W \subset V$ of dimension $n + 1$, then $d_n(B_W)_V = 1$.

Proof : trivial if V is a Hilbert space. Follows from Borsuk-Ulam antipodality theorem in the Banach space case : for any continuous application F from an n -sphere $S_n = \partial B_W$ to an n dimensional space V_n , there exists $x \in S_n$ such that $F(x) = F(-x)$.

It follows that $d_n(\mathcal{K})_V \geq r$ if \mathcal{K} contains the rescaled ball rB_W of an $n + 1$ -dimensional space W . In other words

$$d_n(\mathcal{K})_V \geq b_n(\mathcal{K})_V, \quad n \geq 1,$$

where the Bernstein n -width $b_n(\mathcal{K})_V$ is defined as the largest $r \geq 0$ such that there exists $W \subset V$ of dimension $n + 1$ with $rB_W \subset \mathcal{K}$.

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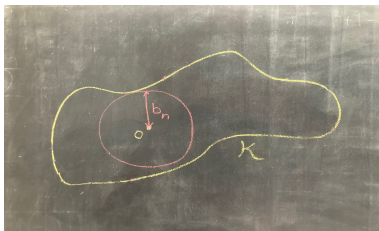
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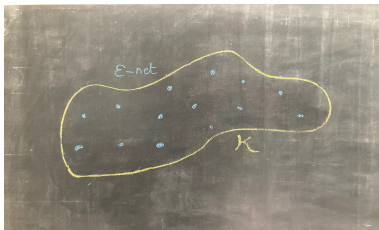
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Entropy numbers

Define $\varepsilon_n(\mathcal{K})_V$ as the smallest ε such that \mathcal{K} can be covered by 2^n balls of radius ε :

$$\mathcal{K} \subset \bigcup_{i=1, \dots, 2^n} B(u^i, \varepsilon), \quad B(u^i, \varepsilon) := \{u : \|u - u^i\|_V \leq \varepsilon\}.$$



Related to lossy coding : Elements of \mathcal{K} can be encoded with n bits up to precision ε_n .

Carl's inequality : for all $s > 0$ one has

$$(n+1)^s \varepsilon_n(\mathcal{K})_V \leq C_s \sup_{m=0, \dots, n} (m+1)^s d_m(\mathcal{K})_V, \quad n \geq 0$$

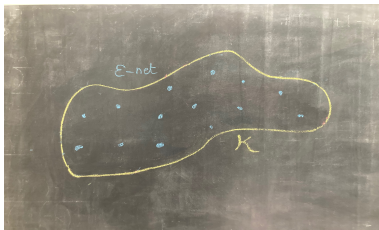
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$$d_n(\mathcal{K})_V \lesssim n^{-s}, \quad n \geq 0 \implies \varepsilon_n(\mathcal{K})_V \lesssim n^{-s}, \quad n \geq 0.$$

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Reduced modeling for parametrized PDEs

Complex problems are often modelled by PDEs involving several physical parameters $y = (y^1, \dots, y^d) \in Y \subset \mathbb{R}^d$.

$$\mathcal{P}(u, y) = 0,$$

For each $y \in Y$, we assume well-posedness and therefore existence of a unique solution $u(y) \in V$.

In certain applications (optimization, inverse problems, uncertainty quantification), we may need to solve $y \mapsto u(y)$ for many instances of $y \in Y$: requires computational methods that are uniformly cheap and efficient, uniformly over $y \in Y$.

We are interested in well approximating the solution manifold

$$\mathcal{K} := \{u(y) : y \in Y\} \subset V,$$

which we assume to be compact.

Reduced modeling usually involves two steps :

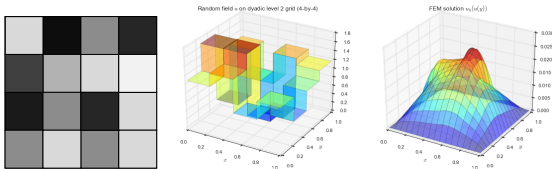
1. In a (**costly**) offline stage, we search for spaces V_n of dimension n that approximate as best as possible the set \mathcal{K} (benchmark $d_n(\mathcal{K})_V$). These spaces are quite different from classical finite element spaces.
2. In a (**cheap**) online stage, for any required $y \in Y$ we may compute an accurate approximation $u_n(y) \in V_n$ of $u(y)$, for example by the Galerkin method.

Estimating n -width of solution manifolds

An instructive example : consider the steady-state diffusion equation

$$-\operatorname{div}(a\nabla u) = f,$$

on a $2d$ domain Ω (+ boundary conditions), with piecewise constant diffusion function $a = a(y)$ having value $\bar{a} + y_j$ on subdomain Ω_j , where $y = (y_1, \dots, y_d) \in Y = [-c, c]^d$.



How large is the n -width of $\mathcal{K} = \{u(y) : y \in Y\} \subset V = H^1(\Omega)$?

Solutions $u(y)$ are bounded in H^s iff $s < 3/2$ and $d_n(\mathcal{U}(H^s))_{H^1} \sim n^{-(s-1)/2} \gtrsim n^{-1/4}$.

In fact $d_n(\mathcal{K})_{H^1}$ decreases faster than $\mathcal{O}(\exp(-cn^{1/d}))$: approximate by power series

$$\max_{y \in Y} \left\| u(y) - \sum_{|\mathbf{v}| \leq k} u_{\mathbf{v}} y^{\mathbf{v}} \right\|_{H^1} \leq C \exp(-ck), \quad y^{\mathbf{v}} = y_1^{v_1} \dots y_d^{v_d},$$

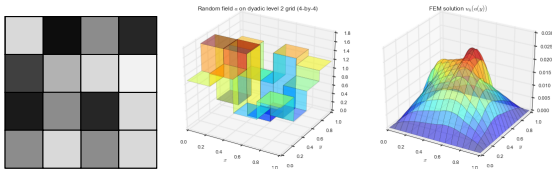
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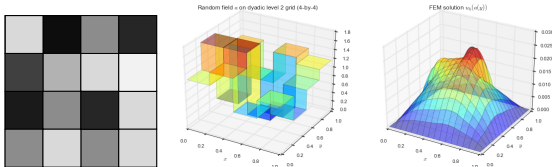
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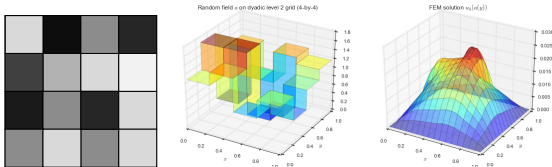
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A general result for infinite dimensional parameter dependence

Theorem (Cohen-DeVore, 2016) : Let V_1 and V_2 be two complex valued Banach spaces and $\mathcal{K}_1 \subset V_1$ be a compact set. Let

$$F : V_1 \rightarrow V_2,$$

be a map that is **holomorphic** on an open neighbourhood of \mathcal{K}_1 . Then, with $\mathcal{K}_2 := F(\mathcal{K}_1)$, one has for all $s > 1$

$$\sup_{n \geq 0} n^s d_n(\mathcal{K}_1)_{V_1} < \infty \implies \sup_{n \geq 0} n^t d_n(\mathcal{K}_2)_{V_2} < \infty, \quad t < s - 1.$$

Note that if F was a continuous linear map, one would simply have

$$d_n(\mathcal{K}_2)_{V_2} \leq C d_n(\mathcal{K}_1)_{V_1}, \quad C = \|F\|_{V_1 \rightarrow V_2}.$$

The proof goes by expanding $a \in \mathcal{K}_1$ in a suitable basis $a = a(y) = \sum_{j \geq 1} y_j \psi_j$ with decay properties on the $\|\psi_j\|_{V_1}$ and then approximate $F(a(y))$ by polynomials in y . This induces a loss of 1 in the rate of decay. **Open problem** : same rate $t = s$?

This result applies to elliptic equations such as $-\operatorname{div}(a \nabla u) = f$ for the map $F : a \rightarrow u$ with $V_1 = L^\infty$ and $V_2 = H^1$. Also applies to parabolic equations, nonlinear problems such as Navier-Stokes equations, and to these problems set on parametrized domains. It does **not** apply to hyperbolic equations.

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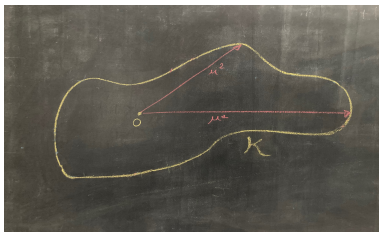
The reduced basis algorithm

Idea : use particular instances $u^i = u(y^i) \in \mathcal{K}$ for generating $V_n = \text{span}\{u^1, \dots, u^n\}$.

Greedy selection in offline stage : having generated u^1, \dots, u^{k-1} , select next instance

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where $P_{V_{k-1}}$ is the orthogonal projection. Here we assume V to be a Hilbert space.



In practice, weak selection $\|u - P_{V_{k-1}} u^k\|_V \geq \gamma \max_{u \in \mathcal{K}} \|u - P_{V_{k-1}} u\|_V$, for fixed $\gamma \in]0, 1[$, and maximization on a large finite training set $\tilde{\mathcal{K}} \subset \mathcal{K}$.

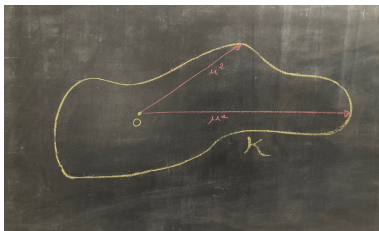
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Approximation performances

For the greedily generated spaces V_n , we would like to compare

$$\sigma_n(\mathcal{K})_V = \text{dist}(\mathcal{K}, V_n)_V = \max_{u \in \mathcal{K}} \|u - P_{V_n} u\|_V,$$

with the n -widths $d_n(\mathcal{K})_V$ that correspond to the optimal spaces.

Direct comparison is deceiving.

Buffa-Maday-Patera-Turinici (2010) : $\sigma_n \leq n 2^n d_n$.

For all $n \geq 0$ and $\varepsilon > 0$, there exists \mathcal{K} such that $\sigma_n(\mathcal{K})_V \geq (1 - \varepsilon) 2^n d_n(\mathcal{K})_V$.

Comparison is much more favorable in terms of convergence rate.

Theorem (Binev-Cohen-Dahmen-DeVore-Petrova-Wojtaszczyk, 2013) : For any $s > 0$,

$$\sup_{n \geq 1} n^s d_n(\mathcal{K})_V < \infty \Rightarrow \sup_{n \geq 1} n^s \sigma_n(\mathcal{K})_V < \infty,$$

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Failure of linear reduced modeling

Linear reduced modeling for parametrized hyperbolic PDEs suffers from a slow decay of Kolmogorov n -width.

Simple example : consider the univariate linear transport equation

$$\partial_t u + a \partial_x u = 0,$$

with constant velocity $a \in \mathbb{R}$ and initial condition $u_0 = u(x, 0) = \chi_{[0,1]}(x)$.

Parametrize the solution by the velocity $a \in [a_{\min}, a_{\max}]$ and consider the solution manifold at final time $T = 1$,

$$\mathcal{H} = \{\chi_{[a, a+1]} : a \in [a_{\min}, a_{\max}]\}.$$

It can be proved that for $1 \leq p < \infty$,

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In particular, we cannot hope for a good performance of reduced basis methods (not better than piecewise constant approximation on uniform meshes).

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Nonlinear approximation

For such problems, one expects improved performance by nonlinear methods.

Non-linear approximation : the function u is approximated by simpler function $v \in \Sigma_n$ that can be described by $\mathcal{O}(n)$ parameters, however Σ_n is not a linear space.

- Rational fractions : $\Sigma_n = \left\{ \frac{p}{q} ; p, q \in \mathbb{P}_n \right\}$.
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where $A_j : \mathbb{R}^{d_j} \rightarrow \mathbb{R}^{d_{j+1}}$ is affine and σ is a nonlinear (rectifier) function applied componentwise, for example $\sigma(x) = \text{RELU}(x) = \max\{x, 0\}$. Here Σ_n is the set of such functions when the total number of parameters does not exceed n .

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Library widths

A library \mathcal{L}_n is a finite collection of linear spaces $V_n \subset V$ of dimension at most n .

We approximate u by picking a space from \mathcal{L}_n , resulting in the error

$$e(u, \mathcal{L}_n)_V = \min_{V_n \in \mathcal{L}_n} \min_{v \in V_n} \|u - v\|_V.$$

Temlyakov (1998) defines the library width

$$d_{N,n}(\mathcal{K})_V := \inf_{\#\mathcal{L}_n \leq N} \max_{u \in \mathcal{K}} e(u, \mathcal{L}_n)_V.$$

Note that $d_{1,n} = d_n$.

The interesting regime is when $N \gg n$. Typical choices that have been studied are $N = A^n$ or $N = n^{an}$ for some $A > 1$ or $a > 0$.

Remark : optimal library approximation amounts in splitting the set \mathcal{K} into N different component \mathcal{K}_j , each of them being approximated by some optimal n -dimensional space V_n^j picked from the library.

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This would lead to the quantity

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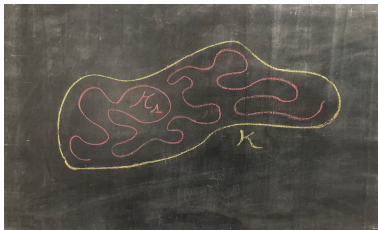
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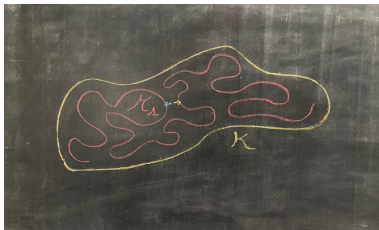
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Both library and manifold widths match known rates of nonlinear approximation (DeVore-Popov, 1980-1990's) by wavelets or adaptive finite elements : if $V = L^p(\Omega)$ and $\mathcal{K} = \mathcal{U}(B_{q,q}^s(\Omega))$ for $\frac{1}{q} < \frac{1}{p} + \frac{s}{d}$, one has

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Upper bounds obtained by these classical nonlinear approximation results.

Library widths satisfy Carl's inequality (for the regimes $N = A^n$ or $N = n^{an}$).

$$(n+1)^s \varepsilon_n(\mathcal{K})_V \leq C_s \sup_{m=0,\dots,n} (m+1)^s d_{m,N}(\mathcal{K})_V, \quad n \geq 0.$$

Manifold widths do **not** satisfy Carl's inequality but are bounded by below by Bernstein widths (by the Borsuk-Ulam argument).

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$$\|D(x) - D(y)\|_V \leq L\|x - y\|_n \quad \text{and} \quad \|E(u) - E(v)\|_n \leq L\|u - v\|_V, \quad x, y \in \mathbb{R}^n, u, v \in V.$$

Here $\|\cdot\|_n$ is an arbitrary norm on \mathbb{R}^n .

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Stable widths and entropies

When V is a Hilbert space, stable widths are strongly tied to entropy numbers.

Theorem : Let V be a Hilbert space, then for any $L > 1$, there exists a constant $c = c(L)$ such that, for any compact set \mathcal{K} ,

$$\delta_{cn,L}(\mathcal{K})_V \leq 3\varepsilon_n(\mathcal{K})_V.$$

With $L = 2$ one can take $c = 26$.

Together with Carl's inequality, this means that

$$\sup_{n \geq 0} n^s \delta_{n,L}(\mathcal{K})_V < \infty \iff \sup_{n \geq 0} n^s \varepsilon_n(\mathcal{K})_V < \infty,$$

for all $s > 0$.

We do not know if this result holds for Banach spaces. Proof for Hilbert spaces :

1. Consider \mathcal{N} an ε_n -net of \mathcal{K} with $\#(\mathcal{N}) = 2^n$.
2. Johnson-Lindenstrauss projection as encoder : $E = P_W$ where $\dim(W) \leq cn$

$$L^{-1} \|u^i - u^j\|_V \leq \|P_W(u^i - u^j)\|_V \leq \|u^i - u^j\|_V, \quad u^i, u^j \in \mathcal{N}.$$

3. This gives an exact decoding map that is L -Lipschitz from $P_W \mathcal{N}$ to \mathcal{N} .
4. Extend this map from $W \sim \mathbb{R}^{cn}$ to V with same Lipschitz constant (Kirschbraun).

Stable widths and entropies

When V is a Hilbert space, stable widths are strongly tied to entropy numbers.

Theorem : Let V be a Hilbert space, then for any $L > 1$, there exists a constant $c = c(L)$ such that, for any compact set \mathcal{K} ,

$$\delta_{cn,L}(\mathcal{K})_V \leq 3\varepsilon_n(\mathcal{K})_V.$$

With $L = 2$ one can take $c = 26$.

Together with Carl's inequality, this means that

$$\sup_{n \geq 0} n^s \delta_{n,L}(\mathcal{K})_V < \infty \iff \sup_{n \geq 0} n^s \varepsilon_n(\mathcal{K})_V < \infty,$$

for all $s > 0$.

We do not know if this result holds for Banach spaces. Proof for Hilbert spaces :

1. Consider \mathcal{N} an ε_n -net of \mathcal{K} with $\#(\mathcal{N}) = 2^n$.
2. Johnson-Lindenstrauss projection as encoder : $E = P_W$ where $\dim(W) \leq cn$

$$L^{-1} \|u^i - u^j\|_V \leq \|P_W(u^i - u^j)\|_V \leq \|u^i - u^j\|_V, \quad u^i, u^j \in \mathcal{N}.$$

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Stable width of solution manifolds

For the linear transport equation manifold $\mathcal{H} = \left\{ \chi_{[a, a+1]} : a \in [a_{\min}, a_{\max}] \right\}$ it is easily established that entropy numbers in L^p spaces have exponential decay

$$\varepsilon_n(\mathcal{H})_{L^p} \leq C \exp(-cn), \quad n \geq 0.$$

This implies in particular that $\delta_{n,L}(\mathcal{H})_{L^2} \leq \tilde{C} \exp(-\tilde{c}n)$ while $d_n(\mathcal{H})_{L^2} \sim n^{-1/2}$.

Similar results hold for manifolds resulting from more general hyperbolic equations.

A general result : if $F : V_1 \rightarrow V_2$ is a L -Lipschitz mapping between Banach spaces, then an ε -net of $\mathcal{K}_1 \subset V_1$ is mapped into an $L\varepsilon$ -net of $\mathcal{K}_2 := F(\mathcal{K}_1)$ and therefore

$$\varepsilon_n(\mathcal{K}_2)_{V_2} \leq L\varepsilon_n(\mathcal{K}_1)_{V_1}, \quad n \geq 0.$$

This implies in particular that when V_2 is a Hilbert space

$$\sup_{n \geq 0} n^s \delta_{n,L}(\mathcal{K}_1)_{V_1} < \infty \implies \sup_{n \geq 0} n^s \delta_{n,L}(\mathcal{K}_2)_{V_2} < \infty.$$

Benchmark : develop concrete stable numerical methods that meet these rates.

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Recent results in this direction

Consider a dictionary \mathcal{D} in a Hilbert space V , such that $\|\varphi\|_V \leq 1$ for all $\varphi \in \mathcal{D}$, and the model class

$$\mathcal{K} := \left\{ v = \sum c_j \varphi_j : \varphi_j \in \mathcal{D}, \sum |c_j| \leq 1 \right\}.$$

The OMP algorithm recursively produces

$$u_n = \sum_{j=1}^n c_j \varphi_j = P_{V_n} u, \quad V_n = \text{span}\{\varphi_1, \dots, \varphi_n\},$$

by selecting $\varphi_n \in \mathcal{D}$ maximizing $\frac{\langle u - P_{V_{n-1}} u, \varphi \rangle}{\|\varphi\|_V}$ over $\varphi \in \mathcal{D}$.

DeVore-Temlyakov (1998) : $u \in \mathcal{K} \implies \|u - u_n\|_V \leq n^{-1/2}$.

Siegel-Xu (2021) : consider smoothly parametrized dictionaries

$$\mathcal{D} = \{\varphi = \mathcal{P}(\theta) : \theta \in \mathcal{M}\},$$

where \mathcal{M} is d -dimensional compact manifold and \mathcal{P} is \mathcal{C}^r . In this setting

$$u \in \mathcal{K} \implies \|u - u_n\|_V \lesssim n^{-s}, \quad s = \frac{1}{2} + \frac{r}{d} \quad \text{and} \quad \varepsilon_n(\mathcal{K})_V \lesssim n^{-s}.$$

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Linear and nonlinear widths measure the approximability of a class \mathcal{K} by linear or nonlinear families of **complexity n** (data compression).

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Are sampling numbers controlled by widths ?

Positive results for **linear** widths : let $V = L^2(D, \mu)$ for some measure μ .

Cohen-Dolbeault (2021) : $\rho_{C_0 n}^{\text{rand}}(\mathcal{K})_V \leq C_1 d_n(\mathcal{K})_V$ for some fixed $C_0, C_1 > 1$.

Temlyakov (2020) : if $\mu(D) < \infty$, then $\rho_{C_0 n}^{\text{det}}(\mathcal{K})_V \leq C_1 d_n(\mathcal{K})_{L^\infty}$ with $C_0 = 1 + \varepsilon$.

These results use weighted least squares as reconstruction and point sparsification techniques due to Spielman, Markus, Srivastava, Nitzan, Olevskii, Ulanovskii.

Negative results for nonlinear widths : there exists classes \mathcal{K} such that sampling numbers $r_n^{\text{det}}(\mathcal{K})_V$ and $r_n^{\text{rand}}(\mathcal{K})_V$ decay much slower than quantities $\delta_{n,L}(\mathcal{K})_V$ or $\varepsilon_n(\mathcal{K})_V$ measuring nonlinear approximation capabilities.

Example : with $V = L^2([0, 1])$, consider the class $\mathcal{K} := \{X_{[a,b]} : a, b \in [0, 1]\}$.

Then $r_n^{\text{det}}(\mathcal{K})_V \sim r_n^{\text{rand}}(\mathcal{K})_V \sim d_n(\mathcal{K})_V \sim n^{-1/2}$, but $\delta_{n,L}(\mathcal{K})_V \sim \varepsilon_n(\mathcal{K})_V \sim \exp(-cn)$.

On the other hand, consider the class $\mathcal{K} := \{X_{[0,a]} : a \in [0, 1]\}$.

Then $d_n(\mathcal{K})_V \sim n^{-1/2}$ but $r_n^{\text{det}}(\mathcal{K})_V \sim \exp(-cn)$ (use adaptive dichotomy).

Question : identify relevant classes of functions for which sampling number behave much better than linear widths, closer to nonlinear widths and entropy numbers.

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Positive results for **linear** widths : let $V = L^2(D, \mu)$ for some measure μ .

Cohen-Dolbeault (2021) : $\rho_{C_0 n}^{\text{rand}}(\mathcal{K})_V \leq C_1 d_n(\mathcal{K})_V$ for some fixed $C_0, C_1 > 1$.

Temlyakov (2020) : if $\mu(D) < \infty$, then $\rho_{C_0 n}^{\text{det}}(\mathcal{K})_V \leq C_1 d_n(\mathcal{K})_{L^\infty}$ with $C_0 = 1 + \varepsilon$.

These results use weighted least squares as reconstruction and point sparsification techniques due to Spielman, Markus, Srivastava, Nitzan, Olevskii, Ulanovskii.

Negative results for nonlinear widths : there exists classes \mathcal{K} such that sampling numbers $r_n^{\text{det}}(\mathcal{K})_V$ and $r_n^{\text{rand}}(\mathcal{K})_V$ decay much slower than quantities $\delta_{n,L}(\mathcal{K})_V$ or $\varepsilon_n(\mathcal{K})_V$ measuring nonlinear approximation capabilities.

Example : with $V = L^2([0, 1])$, consider the class $\mathcal{K} := \{X_{[a,b]} : a, b \in [0, 1]\}$.

Then $r_n^{\text{det}}(\mathcal{K})_V \sim r_n^{\text{rand}}(\mathcal{K})_V \sim d_n(\mathcal{K})_V \sim n^{-1/2}$, but $\delta_{n,L}(\mathcal{K})_V \sim \varepsilon_n(\mathcal{K})_V \sim \exp(-cn)$.

On the other hand, consider the class $\mathcal{K} := \{X_{[0,a]} : a \in [0, 1]\}$.

Then $d_n(\mathcal{K})_V \sim n^{-1/2}$ but $r_n^{\text{det}}(\mathcal{K})_V \sim \exp(-cn)$ (use adaptive dichotomy).

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